

Computer Programs as Dialogue Games

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Meeting of Minds

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Computer Crash

- ▶ On June 4th, 1996 the Ariane 5 expendable launch system was launched into space.
- ▶ 37 seconds later the rocket veered off its flight path and was destroyed by its automated self-destruct system.
- ▶ This was caused by the control software trying to fit a 64-bit number into a 16-bit piece of memory.
- ▶ Sometimes a bug is more than just a nuisance.

Quality Control

- ▶ Computers are everywhere, including many safety-critical situations!
- ▶ So it's important to avoid “bugs”
- ▶ Quality control exists — but prone to complacency / human error
- ▶ Better if we can formally and mechanically check programs.

Formal Methods

- ▶ We study computer programs as objects in their own right.
- ▶ We use
 - ▶ *Mathematics* to describe our programs, how they behave, and what they mean.
 - ▶ *Computer science* to motivate our “design” choices, and implement/use the resulting results

Types and Terms

- ▶ We describe what computer programs look like by using *formal grammars*
 - ▶ $\vdash 3 + 3 : \text{Num}$ says that $3 + 3$ is a valid computer program of type “number”
 - ▶ $\vdash x := x + 1 : \text{Com}$ says that the program that adds 1 to the value of x has type “command”
 - ▶ $\vdash _ + _ : \text{Num} \times \text{Num} \rightarrow \text{Num}$ says that $+$ is a valid program takes two numbers as input and outputs a number.

- ▶ We give *operational semantics* to these well-typed terms to describe how programs *behave*
 - ▶ $3 + 3 \Rightarrow 6$ tells us that the program “3 + 3” reduces to final answer “6”
 - ▶ $(x := x + 1, \{x \mapsto 6\}) \Rightarrow \{x \mapsto 7\}$

Formal Reasoning

- ▶ We now wish to reason about these programs:
 - ▶ Does this program meet its specification?
 - ▶ Is the behaviour of (obviously correct) Program A equivalent to the behaviour of the (more efficient but less clear) Program B?
 - ▶ Is it possible that Program A will go into an infinite loop and get stuck?

Program Equivalence

- ▶ We need to formalise this notion of *equivalence* of two programs
- ▶ Not enough to simply talk about “returning the same answer”
— makes sense for Num, but what about $\text{Num} \rightarrow \text{Com}$?

We say that M_1 for M_2 are *equivalent* iff

Replacing M_1 by M_2 in any program of type Num yields the same answer (and vice-versa).

A Large Quantification

- ▶ Reasoning about program equivalence directly is complex
 - ▶ We have a large quantification — “in **any** program”
- ▶ So we seek *models* of the language that identify equivalent programs

Game Semantics

- ▶ We can model programs as *dialogue games* between the program and its environment.
- ▶ **Program** and **Environment** alternately play *moves*, which might represent an input request or the final answer.
- ▶ Programs are represented as *strategies* for **Program** — recipes that express which move the program should play based on what's happened so far.

Example — Addition

For example, we might consider a program for addition which takes two numbers as input and returns a number ($+ : \text{Num} \times \text{Num} \rightarrow \text{Num}$).

	Num	\times	Num	\rightarrow	Num	
					q	E
	q					P
	n					E
			q			P
			m			E
					$m + n$	P

Game Semantics 1

- ▶ We represent types as games, and terms as strategies.
- ▶ These models are *fully abstract* — equivalent programs are represented by the same strategy.
- ▶ Programming language features correspond to constraints on strategies

⇒ A flexible and accurate way to model programs.

Game Semantics 2

The main reasons this works:

- ▶ **Compositional reasoning** — the meaning of a program is given by the meaning of its components.
- ▶ **Operational, sequential content** — the models are sequential in nature, giving fine grained control of *when* things happen.
- ▶ **Flexibility of constraints** — the models are rich in structure, leading to flexibility for modelling different programming languages

Verification

- ▶ We can check if two programs are correct by seeing if the corresponding strategies are the same.
- ▶ Elsewhere¹, work analyses how this can be done mechanically (and for which languages)
- ▶ We can further use these ideas to check if programs satisfy specifications (transforming program properties to properties on strategies)

¹Ong, Ghica...

Current Work in Bath

Includes:

- ▶ Giving a higher-level, algebraic account of many of the characteristics of game semantics.
- ▶ Investigating the use of game semantics for analysing *access control* and *interference*

Questions?